

Enumeration of certain words *

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1. Introduction

Both the content and methods of this paper are closely related to those of [2]. Let F denote a free semigroup on generators x_1, x_2, \dots, x_m . We wish to enumerate, for fixed n and k , the number of length n words which do not have any segment which is the square of a length k word.

One reason for considering this problem is that it is related to the more difficult problem of enumerating words not containing any segment which is a square word. GOULDEN and JACKSON [1, Corollary 4.1] have obtained our Theorem 1 by completely different methods.

2. A recursion formula

Definition 1. Let F^1 be a free monoid. Let $w_1, w_2 \in F^1$.

- (1) $w_1 | w_2$ if $xw_1y = w_2$ for some $x, y \in F^1$.
- (2) $w_1 |_i w_2$ if $w_2 = w_1y$ for some $y \in F^1$.
- (3) $w_1 |_f w_2$ if $w_2 = xw_1$ for some $x \in F^1$.
- (4) $|w_1|$ denotes the length of w_1 ; $|1| = 0$.
- (5) If $w_1 \neq 1$, \hat{w}_1 is the unique word such that

$$\hat{w}_1 |_i w_1, |\hat{w}_1| = |w_1| - 1.$$

Definition 2. Let F be a free semigroup on $m > 1$ generators and F^1 the associated free monoid. Fix $k > 0$.

To any word w in F^1 we assign a length $k-1$ $(0, 1)$ -vector v as follows. The number v_j is 1 if and only if the $n-k+j+1$ letter of w equals the $n-2k+j+1$ letter of w , and $n-2k+j+1 > 0$. We assign an integer $a(w)$ to w by stating that $a(w)$ is the length of the longest terminal sequence of 1 components of v . Let $S(n, c)$ for $0 \leq c \leq k-1$ denote the set of length n words w_1 of F^1 such that $a(w_1) = c$ and if $|w_2| = k$ it is false that $(w_2)^2 | w_1$.

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In other terms, given a word w_1 consider the terminal segment w_3 of length $2k-1$ of w_1 . Then $a(w_1)$ is the length of the longest initial segment of w_3 which equals a final segment of w_3 , or is $k-1$ if this length exceeds $k-1$. The set $S(n, c)$ is the set of words w of length n such that $a(w)=c$ and w is not divisible by the square of a word of length k .

Theorem 1. For $n > k, 0 < c < k$,

$$|S(n, c)| = |S(n-1, c-1)|,$$

$$|S(n, 0)| = \sum_{j=0}^{k-1} (m-1)|S(n-1, j)|.$$

Furthermore, $|S(n, 0)|=m^n, |S(n, c)|=0$ for $n \leq k, 0 < c < k$.

Proof. Suppose $x \in S(n, c)$. Then $\hat{x} \in S(n, c-1)$. (Note that if $c=k-1$, the $n-k$ and $n-2k$ letters of x must differ, else x would be divisible by the square of a length k word.) And if $y \in S(n, c-1)$ there is one and only one $x \in S(n, c)$ such that $\hat{x}=y$. Namely the last letter of x must equal the $n-k$ letter of y . Thus $|S(n, c)|=|S(n, c-1)|$. Also the function $x \rightarrow \hat{x}$ maps $S(n, 0)$ into $\bigcup_{j=0}^{k-1} S(n-1, j)$. Each element y of the latter set arises from exactly $(m-1)$ elements of $S(n, 0)$. Namely we can add to y any letter except the $(n-k)$ th. Therefore

$$|S(n, 0)| = (m-1) \sum_{j=0}^{k-1} |S(n-1, j)|.$$

This proves the theorem.

This formula can be recast into a matrix form. Let M be the $k \times k$ matrix

$$\begin{bmatrix} m-1 & 1 & 0 & \dots & 0 \\ m-1 & 0 & 1 & \dots & 0 \\ m-1 & 0 & 0 & \dots & 0 \\ \dots & \dots & \dots & \dots & \dots \\ m-1 & 0 & 0 & \dots & 0 \end{bmatrix}.$$

Let u be the vector $(m^k, 0, 0, \dots, 0)$. Then $|S(n, c)|$ is the $c+1$ component of uM^{n-k} , for $n \geq k$. The characteristic polynomial of M is

$$P(z) = z^k - (m-1)(z^{k-1} + z^{k-2} + \dots + z + 1).$$

Definition 3. For $n < k$ put $f(n)=0$. For $n \geq k$, let $f(n)$ be

$$\sum_{j=0}^{k-1} |S(n, j)|.$$

Theorem 2. The generating function of $f(n)$ is

$$\frac{m^k z^k (1-z^k)}{1-mz+(m-1)z^{k+1}}.$$

Proof. As in [2], the generating function of $f(n)$ must have the form

$$\frac{q(z)z^k}{z^k p\left(\frac{1}{z}\right)}$$

where $q(z)$ is some polynomial of degree at most $k-1$. In degree less than $2k$, this quotient must be

$$z^k m^k (1 + mz + \dots + m^{k-1} z^{k-1}).$$

Therefore $q(z)$ is the portion of

$$m^k (1 - (m-1)z - \dots - (m-1)z^k) (1 + mz + \dots + m^{k-1} z^{k-1})$$

having degrees no more than k . A computation gives the formula above (note that $(1-z)$ can be cancelled from numerator and denominator).

3. Asymptotic values of $f(n)$

Lemma 3. *The equation $P(z)=0$ has a unique positive real root r_k of multiplicity 1. This root exceeds the absolute value of any other root. We have $r_k > r_{k-1}$ and*

$$\lim_{k \rightarrow \infty} r_k = m.$$

Moreover

$$m - \frac{m-1}{m^k} > r_k > m - \frac{m-1}{r_j^k}$$

for $j < k$.

Proof. If $P(z)=0$ then $u = \frac{1}{z}$ satisfies

$$u^{k+1} = \frac{mu-1}{m-1}$$

as does $u=1$. However no straight line can cut the curve $y=x^{k+1}$, $x>0$ in more than two places since $y=x^{k+1}$ is concave upwards. Therefore 1 and $\frac{1}{r_k}$ are the only positive solutions of

$$u^{k+1} = \frac{mu-1}{m-1}.$$

Therefore $P(z)=0$ has only one positive real solution r_k and $r_k > 1$. Differentiation shows that r_k has multiplicity 1. Let z be a root of $P(z)=0$ which is negative or complex. Then

$$z^k = (m-1)(z^{k-1} + z^{k-2} + \dots + 1)$$

implies

$$|z|^k < (m-1)(|z|^{k-1} + |z|^{k-2} + \dots + 1).$$

So $|z| < r_k$ because for $|z| \cong r_k$ we have $P(|z|) > 0$. We have

$$r_k^{k-1} = (m-1) \left(r_k^{k-2} + \dots + 1 + \frac{1}{r_k} \right) > (r_k^{k-2} + \dots + 1)(m-1)$$

and

$$r_{k-1}^{k-1} = (m-1)(r_{k-1}^{k-2} + \dots + 1).$$

This implies $r_k > r_{k-1}$. It also implies that

$$r = \lim_{k \rightarrow \infty} r_k$$

satisfies

$$r = (m-1) \frac{1}{1 - \frac{1}{r}}.$$

Therefore $r = m$.

From

$$z^k = (m-1) \frac{z^k - 1}{z - 1}$$

at $z = r_k$ we have

$$z = 1 + (m-1) \left(1 - \frac{1}{z^k} \right).$$

This implies the last inequality of the lemma. This proves the lemma.

Theorem 4. *The asymptotic value of $f(n)$ is*

$$\frac{m^k(1-u^k)u^{k-n-1}}{m-(m-1)(k+1)u^k}$$

where $u = \frac{1}{r_k}$.

Proof. Expand the generating function in partial fractions. All other terms will be insignificant compared with the term

$$\frac{A}{1-r_k z}.$$

This term can be computed by letting z approach $\frac{1}{r_k}$ in the generating function. This proves the theorem.

Abstract

We study the number of words of length n , in m generators, divisible by the square of a length k word. We find a recursion formula, the generating function, and the asymptotic value of this number.

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